Secure data storage



Security

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Problems (1/3)

- > The classical file system protection is limited
 - Physical protection assumptions
 - · Physical confinement of storage devices
 - Logical protection assumptions
 - · Access control performed by systems managing the devices
 - e.g. operating systems
 - Proper use of ACLs or other authorization mechanisms



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Problems (2/3)

- ▶ There are numerous scenarios where this protection is useless
 - Direct/physical access to storage devices
 - · Mobile computational units
 - · Laptops, PDAs, smartphones
 - · Removable storage devices
 - · Tapes, diskettes, CDs DVDs, memory cards
 - Bypassing of logical access control mechanisms
 - · Unethical access by powerful users (e.g. administrators)
 - · Personification of users



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Problems (3/3)

- > Distributed access raises security issues
 - Trust in (unknown) administration teams
 - Remote authentication of users
 - · Security level provided
 - \cdot i.e. how hard it is to impersonate someone
 - · Integration among clients and servers
 - · Applications, operating system
 - · Interaction model
 - · Sessions vs. requests
 - Entities
 - People vs. machines/systems
 - Secure communications
 - · Confidentiality, integrity



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Solution:

File encryption

- - · Can safely circulate along dangerous networks
 - Can safely be stored in insecure storage devices
 - · Either mobile or administrated by others
- ▶ Problems
 - Data retrieval
 - End-users cannot loose encryption/decryption keys
 - · Illegitimate end-user encryption
 - · Corporate data
 - File sharing
 - · It implies some sort of key sharing
 - Interference with regular storage administration procedures
 - e.g. backups



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Ideal architecture (1/2)

- - At the application level
 - · At the level of OS file caches
 - But tacking into consideration authorization issues
- > Visibility of securely stored data
 - Visual awareness
 - Of what is protected and not protected
 - Automatic setting of encryption attributes
 - · With customization options



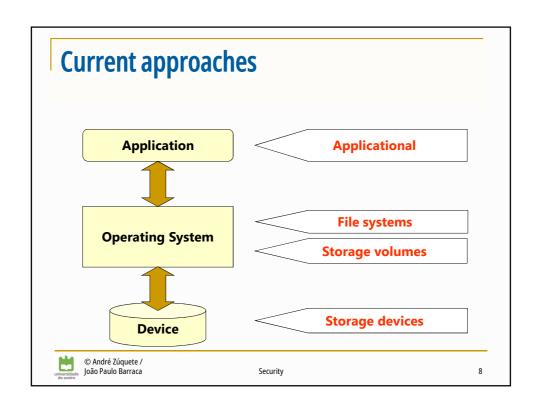
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Ideal architecture (2/2)

- - By groups of users
- Decryption capacity under special circumstances by authorized people
 - Legal enforcement
 - Protection against the loss of decipher keys



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Applicational

- Data transformed by autonomous applications
 - Little or no integration with other applications
 - · Usually it is clear what is secure or not
 - · e.g. using specific file extensions
- - Cleartext resulting files used by other applications
- Data can be transformed with different algorithms
 - · Adds flexibility, increases security
 - Complicates recovery procedures
- - · Secret keys or public keys
- - PGP, AxCrypt, etc.



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Storage volumes / devices

- ▷ Cipher/decipher operations at the volume / device level
 - Total transparency for applications and possibly to the OS
 - The visibility of protected data has volume / device granularity
 - · Not required to handle file systems issues
 - · Protection of meta-information and file data
 - Users and access rights
 - · Cannot differentiate accesses by different users
 - · More suitable for personal storage devices
- - Decipher occurs when data is fetched from devices to server caches
- Examples:
 - · PGPdisk, LUKS (Linux Unified Key Setup)
 - · Self-Encrypting Drives



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Secure file systems: Approaches

- Data is transformed in the path between storage devices and the memory of applications
 - Storage device ⇔ file cache
 - · No protection for remote accesses (server deciphers)
 - · The access to caches gets more complex
 - · Coordination with ACLs
 - · Knowledge of cipher/decipher keys by the OS
 - File cache ⇔ memory of applications
 - · Protection for remote accesses (clients decipher)
 - · Can take place outside the OS (e.g. STDIO in UNIX)
- - · CFS (Cryptographic File System), encfs
 - EFS (Encrypted File System)



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Secure file systems: Limitations (1/2)

- > File system integrity must be preserved
 - · Some file attributes cannot be hidden
 - · For keeping the regular file system operation
 - · Because of other administration tools (e.g. backup tools)
- > Attributes that can easily be hidden
 - Arbitrary file/directory names
 - · Encrypted versions must conform FS naming rules
 - File contents
 - · Preferably without changing file's size



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Secure file systems: Limitations (2/2)

- Attributes that cannot (should not) be hidden/changed
 - Object types
 - · They define the structure of the file system
 - · Contents of directories
 - · Some well-defined names
 - e.g. "." and ".." in UNIX
 - Dates
 - · For managing backups
 - Dimension
 - · For knowing the real occupation of storage devices
 - Ownership
 - · For managing storage quotas
 - Access protection
 - · For keeping the normal access control policies



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Secure file systems: Practical encryption issues

- □ Uniform random access to encrypted data
 - · Ciphers with feedback are not suitable
- - Not advised to use the same key for different files
 - · Similar patterns could reveal similar files
 - Not advised to use the same key for an entire file
 - · Similar patterns along a file could reveal its semantics
 - Stream ciphers are not advised w/ the same key for different files
 - · Known-plaintext attacks could reveal contents of other files



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CFS (Cryptographic File System)

- NFS extension
 - OS ⇔ local CFS server ⇔ local or remote NFS server
 - The NFS interface is kept
 - · The MOUNT interface changes
 - · Includes a password

- · Performed by the local CFS server
 - · Files circulate encrypted in the network
 - · Decrypted file contents are maintained in the client OS file cache
 - All local users with READ access to the file can read the decrypted contents
- · Cipher/decipher keys supplied per each mount point
 - · Communicated to the local CFS server by a modified mount command
 - · This command uses the new MOUNT interface



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CFS

- - Using two keys (K1 and K2) derived from a password
- ▶ Name
 - · Concatenated with and integrity control value
 - Encrypted with ECB
- File contents
 - Stream with OFB ⊕ block ECB
 - · OFB with K1
 - ECB com K2 (disk blocks are not increased)
 - OFB mask computed with K1 per mount point
 - · Random IV per file
 - · Applied between XOR with OFB mask and ECB
 - · Stored in the i-node GID
 - · CFS provides the directory GID instead of the file GID



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EFS (Encrypted File System)

- - First appeared in Windows 2000
 - Provides encryption facilities to NTFS 5
- ▶ Functionality
 - · Each user is bound to an asymmetric key pair
 - · Stored and managed by the OS
 - · Each file is encrypted with a unique symmetric key
 - FEK (File Encryption Key)
 - An encrypted file can be accessed by many users
 - · For each file EFS stores copy of FEK encrypted with the public key of each authorized user
 - Encrypted FEKs are stored in a STREAM associated to the file
 NTFS files are formed by sets of STREAMS
 - Each encrypted file is clearly visible
 - · Using the Explorer file navigator



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EFS cryptographic technology

> Algorithms

- · Asymmetric encryption of FEKs: RSA
- · Symmetric encryption with FEKs: DESX

DESX ≡ DES with whitening

FEK = (K1, K2, K3)

 $C = K1 \oplus DES(K2, P \oplus K3)$

▶ Problems

- · Asymmetric key pairs are stored in disk
 - · Loss risk
 - · Illegitimate access by administrators
- Files are decrypted by servers
 - · No network protection for files stored remotely



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